# Spanning Trees of 2-Complexes from Diagram Groups over the Construction of Semigroup Presentation of Integers using Lifting Method 

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#### Abstract

For any given semigroup presentation we may obtain the fundamental group. In this paper we will determine spanning trees for the 2-complexes of the fundamental groups obtained from the union of two semigroup presentations with finite different initial generators using lifting method. The spanning trees will be systematically selected by using lifting method according to the length of words. Also the general formula for all lifts of spanning trees and the number of edges in the spanning trees will be computed.


Keywords: Fundamental group, semigroup presentation, generators, spanning tree

## 1. INTRODUCTION

The construction for the spanning trees in graphs of semigroup presentations of integers with three and $n$ initial generators namely, ${ }^{3} S=\langle x, y, z \mid x=y, y=z, x=z\rangle$ and ${ }^{n} S=\left\langle x_{1}, x_{2}, x_{3}, \ldots, x_{n} \mid x_{i}=x_{j} ; 1 \leq i<j \leq n\right\rangle$ can be obtained in (see Gheisari and Ghafur (2010) and (2011)) using lifting method. In this research we want to determine spanning tree from fundamental groups over the union of two semigroup presentations of integers with $s$ and $t$ different initial generators by adding a relation.

For given any semigroup presentation $S=\langle X \mid R\rangle$ we may obtain fundamental group $\pi_{1}(K(S))$. Then we can determine the generators from
$\pi_{1}(K(S), U)$ with the basepoint $U$. Thus if $S_{1}=\left\langle X_{1} \mid R_{1}\right\rangle$ and $S_{2}=\left\langle X_{2} \mid R_{2}\right\rangle$, the we compute the $\pi_{1}\left(K\left(S_{1} \cup S_{2}\right), U\right)$.

Guba and Sapir (1997) have shown that if we consider the semigroup presentation $S$, obtained from union of initial generators and relations of two semigroup presentations $S_{1}$ and $S_{2}$ by adding relation $x_{1}=a_{1}$, then $D\left(S, U_{1}\right)$ isomorphic to direct product of $D\left(S_{1}, U_{1}\right)$ and $D\left(S_{2}, U_{2}\right)$. Also they proved in 1997, that if we consider $S=\left\langle X_{1} \cup X_{2} \mid R_{1} \cup R_{2} \cup\left\{U_{1}=U_{2}\right\}\right\rangle$ where $X_{1}, X_{2}$ disjoint sets, and the congruence class of $U_{i}$ modulo $S_{i}$ does not contain words of the form $x U_{i} y$ and $x, y$ are words and $X_{1}, X_{2}$ is not empty. Then $D\left(S, U_{1}\right)$ isomorphic to free product of $D\left(S_{1}, U_{1}\right)$ and $D\left(S_{2}, U_{2}\right)$. Now in this paper, we consider the semigroup presentation $S=\left\langle X_{1} \cup X_{2} \mid R_{1} \cup R_{2} \cup\left\{U_{1}=U_{2}\right\}\right\rangle$ for our method.

Let the two semigroup presentations

$$
{ }^{s} S=\left\langle x_{1}, x_{2}, \ldots, x_{s} \mid x_{i}=x_{j}, 1 \leq i<j\right\rangle
$$

and

$$
{ }^{t} S=\left\langle a_{1}, a_{2}, \ldots a_{t} \mid a_{i}=a_{j}, 1 \leq i<j \leq t\right\rangle
$$

with $s$ and $t$ initial generators. Now we consider the new semigroup presentation

$$
S=\left\langle x_{1}, x_{2}, \ldots, x_{s}, a_{1}, a_{2}, \ldots a_{t} \mid x_{i}=x_{j}, 1 \leq i<j \leq s, a_{i}=a_{j}, 1 \leq i<j \leq t, x_{1}=a_{1}\right\rangle
$$

which is obtained from the union of initial generators and relations of ${ }^{s} S$ and ${ }^{t} S$ by adding a relation $x_{1}=a_{1}$. In this paper we will determine the spanning trees and their lifts of $S$.

In second section we have some preliminaries about diagram groups and semigroup presentation and lifting method. In third section, we will determine the graphs $\Gamma_{n}(S)(n \in N)$ using lifting method. In section results and discussion, we will determine spanning trees of semigroup presentation $S$ according to the length of words, in the graphs $\Gamma_{n}(S)$. Also the general
formula of all lift of spanning trees and the number of edges in spanning trees will be computed.

## 2. PRELIMINARIES

Let $S=\langle X \mid R\rangle$ be a semigroup presentation. Then we may obtain the diagram group $D(S, W)$ where $W$ is a word on $X$ as defined by Guba and Sapir (1997). The 2-complex, associated with presentation $S$ is denoted by $K(S)$. As the 2-complex we may obtain the fundamental group $\pi_{1}(K(S), W)$ with a basepoint $W$. Kilibarda $(1994,1997)$ has shown that the fundamental group $\pi_{1}(K(S), W)$ is isomorphic to diagram group $D(S, W)$. Thus it is sufficient to consider $\pi_{1}(K(S), W)$ instead of $D(S, W)$.

We will consider the fundamental group $\pi_{1}(K(S), W)$ constructed from the semigroup presentation of integers, ${ }^{n} S=\left\langle x_{1}, x_{2}, \ldots, x_{n} \mid x_{i}=x_{j} ; 1 \leq i<j \leq n\right\rangle$. Guba and Sapir (1997) have shown that $\pi_{1}\left(K\left({ }^{3} S\right), x\right)$ is an infinite cyclic, for ${ }^{3} S=\langle x, y, z \mid x=y, y=z, x=z\rangle$.

As the 2-complexes $K(S)$ we may obtain spanning trees of graphs ${ }^{n} \Gamma_{m}(S)$ depending on the length of words. Then we determine the mapping between ${ }^{n} \Gamma_{m}(S)$ and ${ }^{n} \Gamma_{m+1}(S)$. Once we found for ${ }^{n} \Gamma_{1}(S)$, the rest of the graphs are just the lift of ${ }^{n} \Gamma_{1}(S)$.

We will also show that the 2-complex $K(S)$ obtained from semigroup presentation $S$ is actually a union of the graphs ${ }^{n} \Gamma_{m}(S)$ where ${ }^{n} \Gamma_{m}(S)$ contains all vertices of length $m$. Here $n$ refers to the number of initial generators $x_{1}, x_{2}, \ldots, x_{n}$, in the semigroup presentation of integers ${ }^{n} S=\left\langle x_{1}, x_{2}, \ldots, x_{n} \mid x_{i}=x_{j} ; 1 \leq i<j \leq n\right\rangle$. Note that any 2-complex contains vertices, edges, and 2 -cells. Thus a 2 -complex without 2 -cells is simply a graph.

For the semigroup presentation of integers, the 2-complex consists of infinitely connected component ${ }^{n} \Gamma_{m}(S)$ for all $m, n \in N$, where $N$ is a set of the Natural numbers. Note that all vertices in ${ }^{n} \Gamma_{i}(S)$ are words of length $i$. Ahmad and Al-Odhari (2004) proved that if length $(U)=$ length $(V)$ then $\pi_{1}(K(S), U)$ isomorphic to $\pi_{1}(K(S), V)$.

As a group, it is sufficient to determine its generators and relations. The generators of this group can be determined from the 2-complex $K(S)$ by identifying the of a spanning tree $T$. Fix a vertex $v$, where $v$ belong to $K(S)$ and let $e$ be any edge such that $e \notin T$. Then $\gamma_{t(e)} e \gamma^{-1} \tau(e)$ is the generator, where $\gamma_{t(e)}, \gamma_{\tau(e)}$ are paths in a spanning tree $T$ from $v \in K(S)$, to the initial and terminal of $e$ respectively.

Let $U_{i}$ be a word of length $i$. We will show that the generator for $\pi_{1}\left(K(S), U_{i+1}\right)$ can be obtained from the generator of $\pi_{1}\left(K(S), U_{i}\right)$. This is a lifting method. Hence it is sufficient to determine the generator for $\pi_{1}\left(K(S), x_{1}\right)$. Lifting method can determine all generators for the whole groups $\pi_{1}\left(K(S), U_{i}\right)$ for all basepoint $U_{i}$ belongs to $X$. Also using lifting method we can determine the spanning trees of the graphs $\Gamma_{n}(S)$.

## 3. ALGORITHM FOR THE GRAPHS $\Gamma_{n}(S)(n \in N)$

In this section we explain the Algorithm for determining the graphs $\Gamma_{n}(S)$.

Let
$s=\left\langle x_{1}, x_{2}, \ldots, x_{s}, a_{1}, a_{2}, \ldots, a_{t} \mid x_{i}=x_{j}, 1 \leq i<j \leq s, a_{i}=a_{j}, 1 \leq 1<j \leq t, x_{1}=a_{1}\right\rangle$
be a semigroup presentation which is obtained from the union of initial generators and relations of ${ }^{s} S$ and ${ }^{t} S$ by adding a relation $x_{1}=a_{1}$. Associated with semigroup presentation $Q=\langle X \mid R\rangle$ we have a graph $\Gamma$ where the vertices are words on $X$ and the edges are of the form $e=\left(T_{1}, R_{\varepsilon} \rightarrow R_{-\varepsilon}, T_{2}\right)$ such that $t(e)=T_{1} R_{\varepsilon} T_{2}, \quad \tau(e)=T_{1} R_{-\varepsilon} T_{2}$. The graph
obtained from $Q$ is collections of subgraphs $\Gamma_{n}$. Note that the graph $\Gamma\left({ }^{s} S\right)$ obtained from ${ }^{s} S$ is just a collection of subgraphs $\Gamma_{n}\left({ }^{s} S\right)$ where $\Gamma_{n}\left({ }^{s} S\right)$ contains all vertices of length $n$ and respective edges. Similarly we obtain $\Gamma_{n}\left({ }^{t} S\right)$ for ${ }^{t} S$.

Now for $S$, the graph $\Gamma_{n}(S)=\Gamma_{n}\left({ }^{s} S\right) \cup \Gamma_{n}\left({ }^{t} S\right) \cup\left\{\left(u, x_{1} \rightarrow a_{1}, v\right)\right\}$ such that the length $u v=n-1$. If $T_{n}$ is a vertex in $\Gamma_{n}(S)$ then $T_{n} g,\left(g \in\left\{x_{1}, x_{2}, \ldots, x_{s}, a_{1}, a_{2}, \ldots, a_{t}\right\}\right)$ is a vertex in $\Gamma_{n+1}(S)$. Similarly if $\left(u, R_{\varepsilon} \rightarrow R_{-\varepsilon}, v\right)$ is a edge in $\Gamma_{n}(S)$, then $\left(u, R_{\varepsilon} \rightarrow R_{-\varepsilon}, v g\right)$ is the respective edges in $\Gamma_{n+1}(S)$. Thus $\Gamma_{n+1}(S)$ is just $(s+t)$ copies of $\Gamma_{n}(S)$ together with $(s+t)$ vertices $\left(u, x_{1} \rightarrow a_{1}, v g\right) \quad\left(g \in\left\{x_{1}, x_{2}, \ldots, x_{s}, a_{1}, a_{2}, \ldots, a_{t}\right\}\right)$.

For example consider graph $\Gamma_{1}(S)\left(V_{1}, E_{1}\right)$, where $V_{1}=X=\left\{x_{1}, x_{2}, \ldots, x_{s}, a_{1}, a_{2}, \ldots, a_{t}\right\} \quad$ is a set of vertices and $E_{1}=\left\{e_{1 x} \cup e_{1 a} \cup x_{1}=a_{1}\right\} \quad$ is set of edges, where $e_{1 x}=\left\{\left(1, x_{i} \rightarrow x_{j}, 1\right),(1 \leq i<j \leq s)\right\}, e_{1 a}=\left\{\left(1, a_{1} \rightarrow a_{j}, 1\right),(1 \leq i<j \leq t)\right\} \quad$ (see
Figure 1).


Figure 1: Graph of $\Gamma_{1}(S)$

Note that $\Gamma_{2}(S)$ is $(s+t)$ copies of $\Gamma_{1}(S)$ and each vertex in each copy are joined together, respectively by considering the relation $x_{1}=a_{1}$. Similarly, with $(s+t)$ copies of $\Gamma_{2}(S)$, we may obtain $\Gamma_{3}(S)$. Repeating similar procedures for obtain $\Gamma_{4}(S)$ and so on.

Algorithm
Step 1: Determine the graph of $\Gamma_{1}(S)$.
Step 2: The graph $\Gamma_{2}(S)$ is $(s+t)$ copies of $\Gamma_{1}(S)$ similar procedures for obtaining $\Gamma_{n}(S)$ which are $(s+t)$ copies of $\Gamma_{n-1}(S)$.

## 4. RESULTS AND DISCUSSION

In this section we will determine spanning trees in $\Gamma_{n}(S)$. Also the general formula of all lifts of spanning tree and the number of edges in spanning trees will be provided and proved.

## Example 1

Let $T_{1}$ be a spanning tree in $\Gamma_{1}(S)$ where $T_{1}=\left(1, x_{s-1} \rightarrow x_{s}, 1\right)^{-1} \ldots$ $\left(1, x_{1} \rightarrow x_{2}, 1\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, 1\right) \quad\left(1, a_{1} \rightarrow a_{2}, 1\right) \quad\left(1, a_{2} \rightarrow a_{3}, 1\right) \cdots\left(1, a_{t-2} \rightarrow a_{t-1}, 1\right)$ $\left(1, a_{t-1} \rightarrow a_{t}, 1\right)$ (see Figure 2).


Figure 2: Spanning tree in $\Gamma_{1}(S)$
Then the collections all lifts of $\Gamma_{1}$ in $\Gamma_{1}(S)$ at $v_{2}=x_{s} a=\left\{x_{s} x_{1}, x_{s} x_{2}, \ldots, x_{s}^{2}\right.$, $\left.x_{s} a_{1}, \ldots, x_{s} a_{t}\right\}$, for every $a \in X$ are as follows:
(1) Lift of $T_{1}$ at $x_{s} x_{1}$ is:

$$
\begin{gathered}
\left(1, x_{s-1} \rightarrow x_{s}, x_{1}\right)^{-1} \cdots\left(1, x_{1} \rightarrow x_{2}, x_{1}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, x_{1}\right) \\
\left(1, a_{1} \rightarrow a_{2}, x_{1}\right)\left(1, a_{2} \rightarrow a_{3}, x_{1}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, x_{1}\right)\left(1, a_{t-1} \rightarrow a_{t}, x_{1}\right)
\end{gathered}
$$

(2) Lift of $T_{1}$ at $x_{s} x_{2}$ is:

$$
\begin{aligned}
& \left(1, x_{s-1} \rightarrow x_{s}, x_{2}\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, x_{2}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, x_{2}\right) \\
& \left(1, a_{1} \rightarrow a_{2}, x_{2}\right)\left(1, a_{2} \rightarrow a_{3}, x_{2}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, x_{2}\right)\left(1, a_{t-1} \rightarrow a_{t}, x_{2}\right) \\
& \vdots
\end{aligned}
$$

(3) Lift of $T_{1}$ at $x_{s}{ }^{2}$ are:

$$
\begin{aligned}
& \left(1, x_{s-1} \rightarrow x_{s}, x_{s}\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, x_{s}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, x_{s}\right) \\
& \left(1, a_{1} \rightarrow a_{2}, x_{s}\right)\left(1, a_{2} \rightarrow a_{3}, x_{s}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, x_{s}\right)\left(1, a_{t-1} \rightarrow a_{t}, x_{s}\right)
\end{aligned}
$$

and

$$
\begin{aligned}
& \left(x_{s}, x_{s-1} \rightarrow x_{s}, 1\right)^{-1} \ldots\left(x_{s}, x_{1} \rightarrow x_{2}, 1\right)^{-1}\left(x_{s}, x_{1} \rightarrow a_{1}, 1\right) \\
& \left(x_{s}, a_{1} \rightarrow a_{2}, 1\right)\left(x_{s}, a_{2} \rightarrow a_{3}, 1\right) \ldots\left(x_{s}, a_{t-2} \rightarrow a_{t-1}, 1\right)\left(x_{s}, a_{t-1} \rightarrow a_{t}, 1\right)
\end{aligned}
$$

(4) Lift of $T_{1}$ at $x_{s} a_{1}$ is:

$$
\begin{aligned}
& \left(1, x_{s-1} \rightarrow x_{s}, a_{1}\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, a_{1}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, a_{1}\right) \\
& \left(1, a_{1} \rightarrow a_{2}, a_{1}\right)\left(1, a_{2} \rightarrow a_{3}, a_{1}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, a_{1}\right)\left(1, a_{t-1} \rightarrow a_{t}, a_{1}\right)
\end{aligned}
$$

(5) Lift of $T_{1}$ at $x_{s} a_{2}$ is:

$$
\begin{aligned}
& \left(1, x_{s-1} \rightarrow x_{s}, a_{2}\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, a_{2}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, a_{2}\right) \\
& \left(1, a_{1} \rightarrow a_{2}, a_{2}\right)\left(1, a_{2} \rightarrow a_{3}, a_{2}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, a_{2}\right)\left(1, a_{t-1} \rightarrow a_{t}, a_{2}\right) \\
& \vdots
\end{aligned}
$$

(6) Lift of $T_{1}$ at $x_{s} a_{t}$ is:

$$
\begin{aligned}
& \left(1, x_{s-1} \rightarrow x_{s}, a_{t}\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, a_{t}\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, a_{t}\right) \\
& \left(1, a_{1} \rightarrow a_{2}, a_{t}\right)\left(1, a_{2} \rightarrow a_{3}, a_{t}\right) \ldots\left(1, a_{t-2} \rightarrow a_{t-1}, a_{t}\right)\left(1, a_{t-1} \rightarrow a_{t}, a_{t}\right)
\end{aligned}
$$

Example 1 presents all lifts of $T_{1}$ at $v_{1}=x_{s} a, a \in X$, which are exactly a spanning tree in $\Gamma_{2}(S)$.

Theorem 2. Let $T_{n}$ be a collection of all lifts of $T_{1}$ at $x_{1} V_{n-1}$ in $T_{n}(S)$, where $v_{n-1}$ is a word of length $(n-1)$. Then $T_{n}$ is a spanning tree in $T_{n}(S)$.

Proof. By induction on $n$. Consider $T_{2}$ in $\Gamma_{2}(S)$. By definition $T_{2}$ is a collection of lifts and the number of vertices of $T_{2}$ equal to number of vertices in $\Gamma_{2}(S)$, then $T_{2}$ is a spanning tree.

Now suppose $T_{k}$ is a collection of all lifts of $T_{1}$ at $x_{1} V_{k-1}$ in $T_{k}(S)$, thus the number of vertices of $T_{k}$ equal to number of vertices in $\Gamma_{k}(S)$, then $T_{k}$ is a spanning tree. The vertex $x_{1}{ }^{k}$ in the first copy is connected to $x_{2} V_{k-1}, x_{3} V_{k-1}, \ldots, x_{n} V_{k-1}, a_{1} V_{k-1}, a_{2} V_{k-1}, \ldots, a_{t} V_{k-1}$. This is an extra lift of $T_{1}$ at $x_{1} V_{k-1}$ in $\Gamma_{k}(S)$. By definition $T_{k+1}$ is $(s+t)$ copies of $T_{k}$. Similarly $\Gamma_{k+1}(S)$ is $(s+t)$ copies of $\Gamma_{k}(S)$. Hence it is a collection of all lifts of $x_{1} V_{k}$ in $\Gamma_{k+1}(S)$ and the number of vertices of $T_{k+1}$ equal to number of vertices in $\Gamma_{k+1}(S)$. Then $T_{k+1}$ is a spanning tree (see Figure 3).


Figure 3. Spanning tree in $\Gamma_{k}(S)$
Next results show how to compute the total number of lifts in $\Gamma_{n}(S)$ and the number of edges in spanning tree in $\Gamma_{n}(S)$.

Corollary 3. The recurrence formula of all lifts of $T_{n-1}$ in $\Gamma_{n}(S)$ is $l_{n}=(s+t) l_{n-1}+1$ where $l_{i}$ is the total number of lifts of $T_{i},(i=2,3, \ldots)$ in $\Gamma_{i+1}(S)$ and $l_{0}=0$.

Proof. By induction on $n$. For $n=1$ there is only one lift of $T_{1}=\left(1, x_{s-1} \rightarrow x_{s}, 1\right)^{-1} \ldots\left(1, x_{1} \rightarrow x_{2}, 1\right)^{-1}\left(1, x_{1} \rightarrow a_{1}, 1\right)\left(1, a_{1} \rightarrow a_{2}, 1\right)\left(1, a_{2} \rightarrow a_{3}, 1\right)$
$\ldots\left(1, a_{t-2} \rightarrow a_{t-1}, 1\right)\left(1, a_{t-1} \rightarrow a_{t}, 1\right)$ at $v_{1}=x_{s}$ and we denote this number by $l_{1}$. The total number of lifts of $T_{2}$ is $(s+t)+1$, and we denote by $l_{2}$ (refer
to Example 1). Now let $l_{k}$ is the total number of lifts of $T_{k-1}$ in $\Gamma_{k}(S)$ such that $l_{k}=(s+t) l_{k-1}+1$. We will prove that $l_{k+1}$ is the total number of lifts of $T_{k}$ in $\Gamma_{k+1}(S)$ is $l_{k}=(s+t) l_{k-1}+1$. By using the Algorithm $T_{k+1}$ is $(s+t)$ copies of $T_{k}$ plus one (as in proof Theorem 2). Thus, $l_{k+1}=(s+t) l_{k}+1$.

Corollary 4. The total number of lifts of $T_{n-1}$ in $\Gamma_{n}(S)$ is $l_{n}=\frac{(s+t)^{n}-1}{(s+t)-1}$.

Proof. We will prove that by induction. For $n=1$ we have $l_{1}=\frac{(s+t)-1}{(s+t)-1}$.

Then $l_{1}=1$ so its true for $n=1$. Assume true for $n=k$, so $l_{k}=\frac{(s+t)^{k}-1}{(s+t)-1}$. For $n=k+1 \quad$ applying Corollary 3 , we have $l_{k+1}=(s+t) l_{k}+1=(s+t) \cdot \frac{(s+t)^{k}-1}{(s+t)-1}+1=\frac{(s+t)^{k+1}-1}{(s+t)-1}$.

Corollary 5. The recurrence formula of all edges in spanning tree of graph $\Gamma_{n}(S)$ is $e_{n}=(s+t) e_{n-1}+(s+t-1)$, where $e_{n}$ is the total number of edges in spanning tree of $\Gamma_{n}(S)$ and $e_{0}=0$.

Proof. We argue by induction on $n$. For $n=1$, as in Figure 2, the total number of edges in spanning tree of $\Gamma_{1}(S)$ is $(s-1)+(t-1)+1=(s+t-1)$. Now let $e_{k}$ is the total number of edges in spanning tree of $\Gamma_{n}(S)$, that is $e_{k}=(s+t) e_{k-1}+(s+t-1)$ so the formula works when $n=1$. By using the Algorithm and assumption of induction $e_{k+1}$ are $(s+t)$ copies of $e_{k}$ plus $(s+t-1)$ (as in proof Theorem 2 ). Thus, $e_{k+1}=(s+t) e_{k}+(s+t-1)$.

Corollary 6. The total number of edges in the spanning tree $T_{n}$ in $T_{n}(S)$ is $e_{n}=\left((s+t)^{n}-1\right)$.

Proof. By induction on n. For $n=1$ we have $e_{1}=\left((s+t)^{1}-1\right)=(s+t-1)$ (refer to Figure 2). Now let $e_{k}=\left((s+t)^{k}-1\right)$. To prove that $e_{k+1}=\left((s+t)^{k+1}-1\right)$.

By Corollary 5, we conclude

$$
\begin{aligned}
e_{k+1} & \left.=(s+t) e_{k}+(s+t-1)=(s+t) \cdot(s+t)^{k}-1\right)+(s+t-1) \\
& =\left((s+t)^{k+1}-1\right) .
\end{aligned}
$$

## 5. CONCLUSIONS

In this study we determined the new method namely lifting method for finding spanning trees in for 2-complexes of fundamental groups obtained from the union of two semigroup presentations of integers. We also obtained the general formula of all lifts of spanning trees and the number of edges in spanning trees.

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